# Generalized Vertex Induced Connected Subsets of a Graph

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**Abstract**: Let k be a positive integer. A graph G = (V, E) is said to be a  $\Pi_k$  - connected graph if for any given subset S of V(G) with |S| = k, the subgraph induced by S is connected. In this paper, we explore some properties of  $\Pi_k$  - connectedness and its minimality conditions with respect to other graph theoretic parameters.

**Key Words**: Graph, subgraph,  $\Pi_k$  - connected graph, minimal  $\Pi_k$ -connected graph.

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### §1. Introduction

In this article, we consider finite, undirected, simple and connected graphs G=(V,E) with vertex set V and edge set E. As such  $p=\mid V\mid$  and  $q=\mid E\mid$  denote the number of vertices and edges of a graph G, respectively. In general, we use  $\langle X\rangle$  to denote the sub graph induced by the set of vertices  $X\subseteq V$ . N(v) and N[v] denote the open and closed neighborhoods of a vertex v, respectively. A non-trivial graph G is called connected if any two of its vertices are linked by a path in G. A graph G is called n-connected (for  $n\in N$ ) if |V(G)|>n and G-X(the graph that results from removing all vertices in X and all edges incident with these vertices) is connected for any vertex set  $X\subseteq V(G)$  with |X|< n. The greatest integer n such that G is n-connected is called the connectivity  $\kappa(G)$  of G. A cut-edge or cut-vertex of G is an edge or a vertex whose deletion increases the number of components. Unless mentioned otherwise for terminology and notation the reader may refer [1] and [5].

The general problem consists of selecting a set of land parcels for conservation to ensure species availability. This problem is also related to site selection, reserve network design and corridor design. Biologists have highlighted the importance of addressing negative ecological impacts of habitats fragmentation when selecting parcels for conservation. Ways to increase the spatial coherence among the set of parcels selected for conservation have been investigated. Conservation planning via  $\Pi_k$  -connected graph model is an important conservation method, in this model they increase the genetic diversity and allow greater mobility for better response to predation and stochastic events such as fire, as well as long term climate change. This motivated us to study  $\Pi_k$  - connectedness in the following manner:

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For any positive integer k. A graph G is said to be a  $\Pi_k$  - connected graph if for any given subset S of V(G) with |S| = k, the subgraph induced by S is connected.

For more detail, we refer [2]-[4], [6]- [10] and [13].

## §2. $\Pi_k$ -Connected Graphs

**Proposition** 2.1 For any graph G with  $p \ge 1$  vertices is a  $\Pi_1$  - connected graph.

**Proposition** 2.2 For a given  $p \geq 3$  vertices, there exist a  $\Pi_3$  - connected graph of a graph G.

*Proof* Removal of t independent edges from a complete graph on p vertices results into a  $\Pi_3$  - connected graph, where  $0 \le t \le \frac{p}{2}$  if p is even and  $0 \le t \le \frac{p-1}{2}$  if p is odd.

**Proposition** 2.3 Let  $\xi$  be the number of edges required to make a totally disconnected graph which is a  $\Pi_3$  - connected graph and hence

$$\xi(\overline{K}_p) = \begin{cases} \frac{p^2 - 2p}{2} & \text{if } p \text{ is even} \\ \frac{p^2 - 2p + 1}{2} & \text{if } p \text{ is odd} \end{cases}$$

For a complete bipartite graph  $K_{m,n}$ , the number of edges to be added to make it a  $\Pi_3$  - connected graph is given by  $\xi(\overline{K}_m) + \xi(\overline{K}_n)$ .

**Proposition** 2.4 In general the number of edges required to make complete n- partite graph  $K_{x_1,x_2,x_3,...,x_n}$  as a  $\Pi_3$  - connected graph is given by

$$\xi(\overline{K}_{x_1}) + \xi(\overline{K}_{x_2}) + \cdots + \xi(\overline{K}_{x_n}).$$

**Proposition** 2.5 The complete bipartite graph  $K_{m,n}$  is a  $\Pi_3$  - connected graph if m = 1, 2 and n = 1, 2.

**Theorem** 2.1 For any graph G with  $p \geq 3$  vertices is a  $\Pi_3$  - connected graph if and only if  $deg(v_i) \geq p-2$  for all  $v_i \in V(G)$ .

Proof Let G be a  $\Pi_3$  - connected graph. Suppose on contrary, there exists  $v_j$  such that  $deg(v_j) \leq p-3$ . Let  $v_1$  and  $v_2$  be any two vertices which are not adjacent to  $v_j$ , thus the graph induced by  $v_1, v_2$  and  $v_j$  is not connected, which is a contradiction. Hence  $deg(v_i) \geq p-2$  for all  $v_i \in G$ .

Conversely, suppose  $deg(v_i) \ge p-2$  for all  $v_i \in G$ , let  $v_1, v_2, v_3$  be any three vertices. Then  $v_1$  is adjacent to at least one of  $v_2$  and  $v_3$ .  $v_2$  is adjacent to at least one of  $v_1$  and  $v_3$ . Also  $v_3$  is adjacent to at least one of  $v_1$  and  $v_2$ . Therefore G is a  $\Pi_3$  - connected graph.

**Theorem** 2.2 If a graph G is a  $\Pi_k$ - connected graph, then

$$\delta(G) \ge \begin{cases} t & \text{if } p = t(k-1) + 1 \\ t+1 & \text{if } p = t(k-1) + r + 1, 1 \le r \le k - 2 \end{cases}$$

Proof Let G be a  $\Pi_k$  connected graph with p = t(k-1) + 1 vertices. Suppose on contrary that  $\delta(G) = s < t$ . Let v be a vertex with deg(v) = s. Now we partition the remaining p-1 vertices into t vertex disjoint sets such that each set contains a vertex adjacent to v. Since s < t, there exists at least one set N which has no vertex adjacent to v. Then  $\langle N \cup \{v\} \rangle$  is a disconnected subgraph on k vertices, which is a contradiction.

Now, let p-1=t(k-1)+r. Suppose on contrary that  $\delta(G)=s< t+1$ , let v be a vertex such that deg(v)=s. Now partition the remaining p-1 vertices excluding the vertex v in t+1 number of vertex disjoint subsets having t subsets with cardinality k-1 and one with cardinality r in such a way that each subset contains exactly one vertex adjacent to v. Since deg(v)< t+1, there exist at least one subset having no vertex adjacent to v. If the cardinality of such a subset, say T is k-1 then  $\langle T \cup \{v\} \rangle$  is a disconnected subgraph on k- vertices, again a contradiction. If the cardinality of such a subset, say D is r, then take a vertex which is adjacent to v from a subset A with cardinality k-1 to D and any one vertex from D to A, then  $\langle A \cup \{v\} \rangle$  is a disconnected subgraph which is induced by k- vertices, a contradiction. Hence

$$\delta(G) \ge \begin{cases} t & \text{if } p - 1 = t(k - 1) \\ t + 1 & \text{if } p - 1 = t(k - 1) + r, 1 \le r \le k - 2 \end{cases} \square$$

**Theorem** 2.3 Let G be a  $\Pi_k$  - connectedHence the result follows. Then G is a  $\Pi_{k+1}$  - connected graph with  $2 \le k \le p-1$ .

Proof On contrary, suppose G is a  $\Pi_k$  - connected graph but not a  $\Pi_{k+1}$  - connected one. Let S be set of k+1 vertices on which the graph induced has more than one component. Clearly a subset T consisting of k vertices on which the graph induced is disconnected which is a subgraph of G, which is a contradiction. Hence G is also a  $\Pi_{k+1}$  - connected graph with  $2 \le k \le p-1$ .

**Theorem** 2.4 A graph G is a  $\Pi_k$ -connected graph for all k,  $1 \le k \le p$  if and only if G is isomorphic to  $K_p$ .

*Proof* Let G be a  $\Pi_k$  - connected graph for all k,  $1 \le k \le p$ . Clearly G is isomorphic to a complete graph  $K_p$  as G is a  $\Pi_2$  - connected graph.

On the other hand, let G be a complete graph on p vertices. In  $K_p$ , every pair of vertices are adjacent. Hence G is a  $\Pi_2$  - connected graph. As we have proved, every  $\Pi_2$  - connected graph is  $\Pi_t$  - connected graph for all t,  $3 \le t \le p$ . Therefore G is a  $\Pi_k$  - connected graph with  $1 \le k \le p$ .

**Theorem** 2.5 For any Square tree  $T^2$  of a tree T with diameter  $d(T) \geq 5$  is a  $\Pi_{p-1}$  - connected

graph.

Proof Let T be any tree with diameter  $d(T) \geq 5$ . Consider any four non pendent vertices  $v_{i+1}, v_{i+2}, v_{i+3}, v_{i+4}$  in T such that  $v_{i+1}$  is adjacent to  $v_{i+2}, v_{i+2}$  is adjacent to  $v_{i+3}, v_{i+3}$  is adjacent to  $v_{i+4}$ , then in square of T,  $v_{i+2}$  is adjacent to  $v_{i+1}, v_{i+3}, v_{i+4}$  and  $v_{i+3}$  is adjacent to  $v_{i+1}, v_{i+2}, v_{i+4}$ . Removal of  $v_{i+2}$  and  $v_{i+3}$  disconnects the square and since no vertex is a cut vertex, removal of any one vertex does not disconnect  $T^2$ . Hence square of any tree T with diameter  $d(T) \geq 5$  is a  $\Pi_{p-1}$  - connected graph.

**Proposition** 2.6 For any integer  $m \ge 1$ ,  $K_{1,m}$ ,  $K_{1,m}^2$  are  $\Pi_2$  - connected graphs.

*Proof* Let  $K_{1,m}$  be any star. Then the square of a star is a complete graph on m+1 vertices and hence a  $\Pi_2$  - connected graph.

#### §3. Minimality Conditions on $\Pi_k$ -Connected Graphs

A  $\Pi_k$  - connected graph G is said to be a vertex minimal  $\Pi_k$  - connected graph if G is not a  $\Pi_{k-1}$  - connected graph. A vertex minimal  $\Pi_k$  - connected graph G is said to be a partially vertex - edge minimal  $\Pi_k$  - connected graph if G - e is not a  $\Pi_k$  - connected graph for some  $e \in E(G)$ . A vertex minimal  $\Pi_k$  - connected graph G is said to be a totally vertex - edge minimal  $\Pi_k$  - connected graph if G - e is not a  $\Pi_k$  - connected graph for every  $e \in E(G)$ . For more details, refer [12].

**Proposition** 3.1 For any cycle  $C_p$ ;  $p \ge 4$  vertices, the number of edges to be added to make it a  $\Pi_3$  - connected graph is given by

$$q \ge \begin{cases} \frac{p^2 - 4p}{2} & \text{if } p \text{ is even} \\ \frac{p^2 - 4p + 1}{2} & \text{if } p \text{ is odd} \end{cases}$$

where, the equality holds when the resulting graph is a partially vertex - edge minimal  $\Pi_3$  -connected graph.

*Proof* Let  $C_p$ ;  $p \ge 4$  vertices be a cycle. Then we have the following cases.

Case 1. Suppose p is even. In any  $\Pi_3$  - connected graph, the degree each vertex is at least p-2. Hence the number of edges in any  $\Pi_3$  - connected graph is always greater than or equal to  $\frac{p(p-2)}{2}$ . Therefore the number of edges to be added to  $C_p$  is greater than or equal to  $\frac{p^2-2p}{2}-p=\frac{p^2-4p}{2}$ .

Case 2. Suppose p is odd. In any  $\Pi_3$  - connected graph on odd number of vertices, the degree of each of the p-1 vertices is at least p-2 and the degree of one vertex is p-1. Hence the number of edges in any  $\Pi_3$  - connected graph on odd number of vertices is always greater than or equal to  $\frac{(p-1)^2}{2}$ . Hence the number of edges to be added to  $C_p$  is greater than or equal to  $\frac{p^2-2p+1}{2}-p=\frac{p^2-4p+1}{2}$ .

**Proposition** 3.2 For any path  $P_p$ ;  $p \ge 4$  vertices, the number of edges to be added to make it a  $\Pi_3$  - connected graph is given by

$$q \ge \begin{cases} \frac{p^2 - 4p}{2} + 1 & \text{if } p \text{ is even} \\ \frac{p^2 - 4p + 1}{2} + 1 & \text{if } p \text{ is odd} \end{cases}$$

where the equality holds when the resulting graph is a partially vertex - edge minimal  $\Pi_3$  -connected graph.

*Proof* The proof follows on the same lines as in the above proposition.

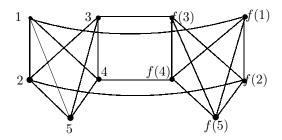
**Theorem** 3.1 A connected graph G is a vertex minimal  $\Pi_p$  - connected graph if and only if it has at least one cut vertex.

Proof Let a connected graph G be a  $\Pi_p$  - connected graph, that is, G is not a  $\Pi_{p-1}$  -connected graph. There exist a vertex v such that the graph induced by V(G)-v is disconnected. Hence v is cut vertex. Conversely, let G be a connected graph with a cut vertex, say v, therefore the subgraph induced by the vertices V(G)-v is disconnected. Hence the graph G is a vertex minimal  $\Pi_p$  - connected graph.

**Theorem** 3.2 For a given k = 2l + 1,  $l \ge 3$ , there exists  $\Pi_k$  - connected graph.

Proof Let  $k=2l+1,\ l\geq 3$  and G be a  $\Pi_3$  - connected graph on k-3 vertices with  $V(G)=\{1,2,3,\cdots,k-3\}$  and let G' be a graph with  $V(G')=\{1,2,3,\cdots,k-3,k-2\}$  obtained by adding a vertex k-2 and making it adjacent to all the vertices of G. Now take prism of G', label the vertices of second copy of G' in the prism as  $\{f(1),f(2),f(3),\cdots,f(k-3),f(k-2)\}$  such that f(i) is the mirror image of i and remove the edge (k-2,f(k-2)) from the prism. In the resulting graph H (say), the subgraph induced by any subset  $S\subseteq V(H)/\{k-2,\ f(k-2)\}$  containing k-1 vertices is connected. The subgraph induced by  $V(G)\cup\{k-2,\ f(k-2)\}$  disconnected on k-3+1+1=k-1 vertices and hence every subgraph induced by k vertices is connected. Hence k is k - connected.

**Observation** 3.1 The graph obtained in the above theorem is regular when G is a partially vertex - edge minimal  $\Pi_3$  - connected graph, which is having even order.



**Figure 1.** Prism of a  $\Pi_3$  - connected graph having odd order

For illustration, we construct the above prism of a  $\Pi_3$  - connected graph having odd order, where prism of a graph G is defined as the cartesian product  $G \times K_2$ ; that is, take two disjoint copies of G and add a matching joining the corresponding vertices in the two copies, [8].

**Theorem** 3.3 For any vertex minimal  $\Pi_k$  - connected graph can be embedded in a vertex minimal  $\Pi_{k+i}$  - connected graph, where  $i \geq 0$ .

Proof Let  $G_1$  and  $G_2$  be two vertex minimal  $\Pi_k$  and  $\Pi_{k+i}$  - connected graphs. Now we construct a vertex minimal  $\Pi_{k+i}$  - connected graph in which  $G_1$  is an induced vertex minimal  $\Pi_k$  - connected subgraph. Make each vertex of  $G_1$  adjacent to each vertex in  $G_2$  and let the resulting graph be G. Let S be any set of k+i vertices from G. In the following cases we prove the graph induced by S is connected.

**Case 1.** Suppose  $S \cap V(G_1) \cap V(G_2) \neq \emptyset$ , then the graph induced by S is connected since each vertex in  $S \cap V(G_1)$  is adjacent to every vertex of  $S \cap V(G_2)$ .

Case 2. Suppose  $S \cap V(G_1) \neq \emptyset$  and  $S \cap V(G_2) = \emptyset$ , then  $|S| \geq k$  and  $S \subseteq V(G_1)$ , the graph induced by S is connected as  $G_1$  is vertex minimal  $\Pi_k$  - connected.

Case 3. Suppose  $S \cap V(G_1) = \emptyset$  and  $S \cap V(G_2) \neq \emptyset$ , then the graph induced by S is connected since S is completely contained in  $V(G_2)$  and  $G_2$  is a vertex minimal  $\Pi_{k+i}$  - connected graph. In all the three cases the graph induced by S is connected and G is not  $\Pi_{k+i-1}$  - connected graph since  $G_2$  is a vertex minimal  $\Pi_{k+i}$  - connected graph. Hence the graph G is a vertex minimal  $\Pi_{k+i}$  - connected graph having vertex minimal  $\Pi_k$  - connected graph  $G_1$  as its induced subgraph.

Thus the result follows.  $\Box$ 

**Theorem** 3.4 A connected graph G is a partially vertex - edge minimal  $\Pi_p$  - connected graph if and only if it has at least one cut edge.

Proof Let G be a connected graph with a cut edge say e = uv. Here u is a cut vertex and also G - e is disconnected, hence G is a partially vertex - edge minimal  $\Pi_p$  - connected graph. Conversely, let G be a partially vertex - edge minimal  $\Pi_p$  - connected graph then G - e is not a  $\Pi_p$  - connected graph. Hence e is a cut edge.

To prove our next result we make use of the following observations.

Observation 3.2 Removal of  $\frac{p}{2}$  independent edges from a complete graph  $K_p$  on even number of vertices results into a partially vertex - edge minimal  $\Pi_3$  - connected graph having  $\frac{p(p-2)}{2}$  edges.

**Observation** 3.3 Removal of  $\frac{p-1}{2}$  independent edges from a complete graph on odd number of vertices results into a partially vertex - edge minimal  $\Pi_3$  - connected graph having  $\frac{(p-1)^2}{2}$  edges.

Observation 3.4 Partially vertex - edge minimal  $\Pi_3$  - connected graph having even order is a regular graph with regularity p-2.

**Theorem** 3.5 Let  $G_1$  and  $G_2$  be two partially vertex - edge minimal  $\Pi_3$  - connected graph. If  $V(G_1) = V(G_2)$ , then  $G_1$  is isomorphic to  $G_2$ .

*Proof* Let  $G_1$  and  $G_2$  be partially vertex - edge minimal  $\Pi_3$  - connected graph having same order. Then, there are following cases:

Case 1. Suppose p is even. As the graphs  $G_1$  and  $G_2$  are partially vertex - edge minimal  $\Pi_3$  connected graphs, deg(v) = p - 2, for all  $v \in V(G_1)$  and deg(v) = p - 2 for all  $v \in V(G_2)$ . Clearly  $G_1$  is isomorphic to  $G_2$ .

Case 2. Suppose p is odd. As the graphs  $G_1$  and  $G_2$  are partially vertex - edge minimal  $\Pi_3$  -connected graphs, in the graphs  $G_1$  and  $G_2$  degree of each of p-1 vertices is p-2 and degree of one vertex is p-1. Hence in this case also  $G_1$  is isomorphic to  $G_2$ .

**Theorem** 3.6 Any graph G with order p is a totally vertex - edge minimal  $\Pi_p$  - connected graph if and only if G is isomorphic to a tree on p vertices.

Proof Let G be a graph of order p which is a strongly critical  $\Pi_p$  - connected graph, that is, G-e is not  $\Pi_p$  - connected for all  $e \in V(G)$  and G is not a  $\Pi_{p-1}$  - connected graph. The first condition in a totally vertex - edge minimal  $\Pi_p$  - connected graph which implies that every edge in G is a bridge and the second condition in a totally vertex - edge minimal  $\Pi_p$  - connected graph which again implies that every vertex in G is a cut vertex, clearly G is isomorphic to a tree on p vertices.

Conversely, let G is isomorphic to a tree on p vertices. Since every internal vertex is a cut vertex, we have a disconnected induced subgraph on p-1 number of vertices and since every edge is a bridge, G-e is not a  $\Pi_p$ - connected graph for all  $e \in V(G)$ . Hence G is a totally vertex - edge minimal  $\Pi_p$  - connected graph.

**Theorem** 3.7 Any graph G having even number of vertices is a totally vertex - edge minimal  $\Pi_3$  - connected graph if and only if  $deg(v_i) = p - 2$  for all  $v_i \in V(G)$ .

Proof Let G be a totally vertex - edge minimal  $\Pi_3$  - connected graph on even number of vertices, implies  $deg(v) \geq p-2$  for all  $v \in V(G)$  from the Theorem 2.1. Suppose deg(v) > p-2 for some  $v \in V(G)$ , i.e., deg(v) = p-1. There exist a vertex w adjacent to v such that deg(w) = p-1. The graph G - vw is still a  $\Pi_3$  - connected graph. Hence G is not a totally vertex - edge minimal  $\Pi_3$  - connected graph, which is a contradiction. Hence deg(v) = p-2 for all  $v \in V(G)$ .

Conversely, let G be a graph such that deg(v) = p - 2 for all  $v \in V(G)$ . For every vertex in G there exist an unique non adjacent vertex in G. Hence in G - e there exist two vertices say v and w non adjacent to some vertex say u. The graph induced by these three vertices is disconnected and hence the graph is a totally vertex - edge minimal  $\Pi_3$  - connected graph.  $\square$ 

**Observation** 3.5 Any complete graph on  $p \geq 3$  vertices is a  $\Pi_3$  - connected graph but not a partially vertex - edge minimal and totally vertex - edge minimal  $\Pi_3$  - connected graph.

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